# Computational Thinking (CO2412): Tutorial - Calendar Week 43 

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### 1.3.1. Final digit enumeration problem

In the lecture, we discussed an iterative algorithm and its Python implementation, called mod10_count_naive () in the associated Jupyter Notebook, ${ }^{1}$ for the following problem:

The input consists of two arguments, a list $\mathbf{x}=\left[x_{0}, x_{1}, \ldots, x_{n-1}\right]$ of $n=\operatorname{len}(\mathbf{x})$ integer numbers, where multiple elements are allowed to have the same value, and a single-digit integer $y$ with $0 \leq y \leq 9$. A list $\left[q_{1}, q_{2}, q_{3}\right]$ is returned where:

1. $q_{1}$ is the number of indices $i$ such that $x_{i}$ has $y$ as its final digit. Differently expressed, it is the number of list elements such that $x_{i} \bmod 10=y$, where mod stands for modulo, i.e., remainder after division. ${ }^{2}$ If the same number occurs multiple times in the list, it also counts multiple times, once for each index.
2. $q_{2}$ is the number of ordered pairs $(i, j)$ of indices, with $i \neq j$, such that $x_{i} x_{j} \bmod 10=y ;$ i.e., $y$ is the final digit of $x_{i} x_{j}$. The two different ways of arranging the indices, $(i, j)$ and $(j, i)$, both count separately - therefore, $q_{2}$ is always an even number. Note that the requirement is for $i$ and $j$ to be different, not $x_{i}$ and $x_{j}$.
3. $q_{3}$ is the number of ordered triples $(i, j, k)$ of indices, all different from each other $(i \neq j, i \neq k, j \neq k)$, for which $x_{i} x_{j} x_{k} \bmod 10=y$. As above, all the different permutations (i.e., arrangements) of the three indices each count separately, of which there are six each time; accordingly, $q_{3}$ is always divisible by 6 .

For example, if $\mathbf{x}=[24,8,19,8,2]$ and $y=4$, the list $[1,2,24]$ needs to be returned. ${ }^{3}$ The mod10_count_naive() code solves this problem, but it has $\mathrm{O}\left(n^{3}\right)$ time requirements, by which it does not perform very favourably for long lists.
a) Propose a more efficient algorithm and develop a more performant code.
b) Of what order is the time efficiency of your algorithm, using Landau notation (i.e., "big O notation")? Provide a brief justification similar to those from the lecture.

[^0]c) Conduct performance measurements, including but not necessarily limited to the two demo lists $\times 200$ and $\times 1000$ from the notebook, ${ }^{4}$ with $n=200$ and 1000, respectively. What is the ratio between the two runtimes? For the naive implementation, which scales with $\mathrm{O}\left(n^{3}\right)$, it is close to $125=(1000 / 200)^{3}$; for a code that has an asymptotic runtime in $\mathrm{O}\left(n^{m}\right)$, a ratio close to $5^{m}$ should be expected.

### 1.3.2. Number matching problem

The function natmatch_iter() takes two arguments: First, a list of $k$ integer numbers $\mathbf{x}=\left[x_{0}, x_{1}, \ldots, x_{k-1}\right]$, and second, a natural number $y$; it determines whether there is a match, here defined by the existence of two list elements with $x_{i}+x_{j}=y$, where $x_{i} \neq x_{j}$.

In the present and the previous notebook, we were calling this function for a given value of $k$ many times, where the $k$ elements of the list $\mathbf{x}$ were assigned new random values each time, using a uniform random distribution ${ }^{5}$ over all integers from 0 to $k^{2}-1$. The second argument was given by $y=k^{2}$. Statistics from these function calls make it apparent that for large values of $k$, a match is found in about $39 \%$ to $40 \%$ of the cases.

Determine the fraction of cases for which there is a match, in the case of large $k$ (ideally, as $k$ approaches infinity), as accurately as possible. ${ }^{6}$

Submission deadline: 13th November 2021; discussion planned for 25th November 2021. Group work by up to four people is welcome.

[^1]
[^0]:    ${ }^{1}$ For the notebook, cf. https://home.bawue.de/ ${ }^{\sim}$ horsch/teaching/co2412/material/ iterative-algorithms.ipynb.
    ${ }^{2}$ In Python, this condition is expressed by x[i] \% $10==\mathrm{y}$.
    ${ }^{3} q_{1}=1$ for $x_{0}=24, q_{2}=2$ for $x_{1} x_{3}=x_{3} x_{1}=64$, and $q_{3}=24$ for $x_{1} x_{2} x_{4}=x_{1} x_{4} x_{2}=x_{2} x_{1} x_{4}=x_{2} x_{3} x_{4}=$ $x_{2} x_{4} x_{1}=x_{2} x_{4} x_{3}=x_{3} x_{2} x_{4}=x_{3} x_{4} x_{2}=x_{4} x_{1} x_{2}=x_{4} x_{2} x_{1}=x_{4} x_{2} x_{3}=x_{4} x_{3} x_{2}=304$, in combination with $x_{0} x_{1} x_{4}=$ $x_{0} x_{3} x_{4}=x_{0} x_{4} x_{1}=x_{0} x_{4} x_{3}=x_{1} x_{0} x_{4}=x_{1} x_{4} x_{0}=x_{3} x_{0} x_{4}=x_{3} x_{4} x_{0}=x_{4} x_{0} x_{1}=x_{4} x_{0} x_{3}=x_{4} x_{1} x_{0}=x_{4} x_{3} x_{0}=384$.

[^1]:    ${ }^{4}$ For validation, the return value for $\mathbf{x}=\mathrm{x} 200, y=7$ should be [28, 1528, 134610], and for $\mathbf{x}=\mathrm{x} 1000$, $y=7$ it should be $[105,42660,17483370]$.
    ${ }^{5}$ That is, each integer from 0 to $k^{2}-1$ had the same probability of being assigned to any of the list elements.
    ${ }^{6}$ The method suggested here is to run a large number of function calls with random input for a large value of $k$, by which a sufficient accuracy should be reached. With some mathematical knowledge, going beyond the scope of this module, is also possible to give an exact answer; note, however, that here you are not expected to do this (of course, any such solutions or attempts are nonetheless very welcome).

