

Cashier problem (T2.2)

The cashier problem is specified as follows. The function solving the problem has two arguments:

- 1) first, a natural number, given in the smallest currency unit (e.g., pence), representing an **amount of money** that is to be paid out;
- 2) second, a sorted list with the **values of the existing coin types**, in the same currency unit (we assume that "1" is always among these values).

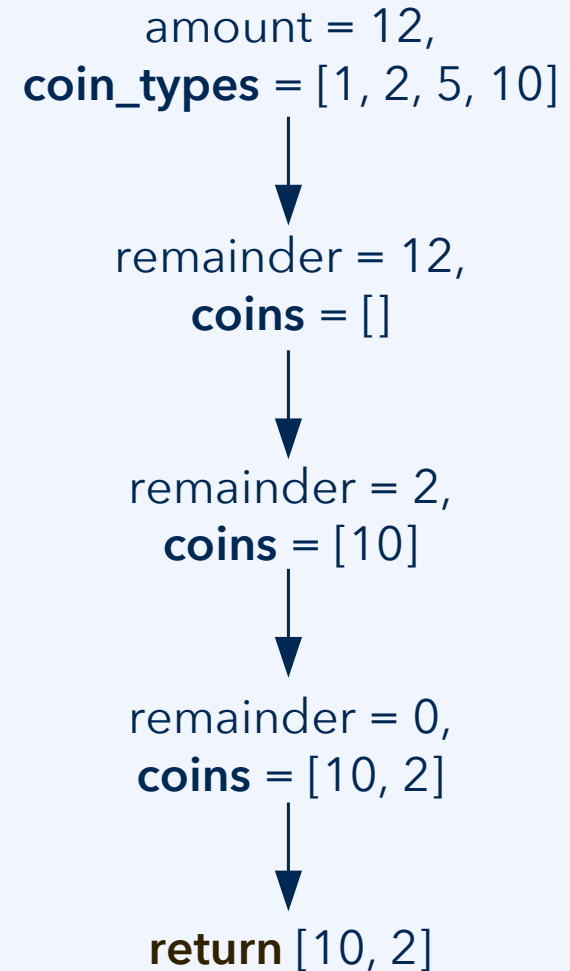
As the function's return value, we expect a **list containing coin values** that add up to the requested amount; this must be the **shortest possible list**, *i.e.*, we want to use as few coins as possible.

Note that as a precondition it is assumed that the list passed as the function's second argument is already sorted.

Cashier problem (T2.2)

greedy algorithm

```
def cashier(amount, coin_types):  
    coins = []  
    remainder = amount  
  
    while remainder >= coin_types[0]:  
  
        for i in range(len(coin_types)-1, -1, -1):  
            if remainder >= coin_types[i]:  
                coins.append(coin_types[i])  
                remainder -= coin_types[i]  
                break  
        return coins
```



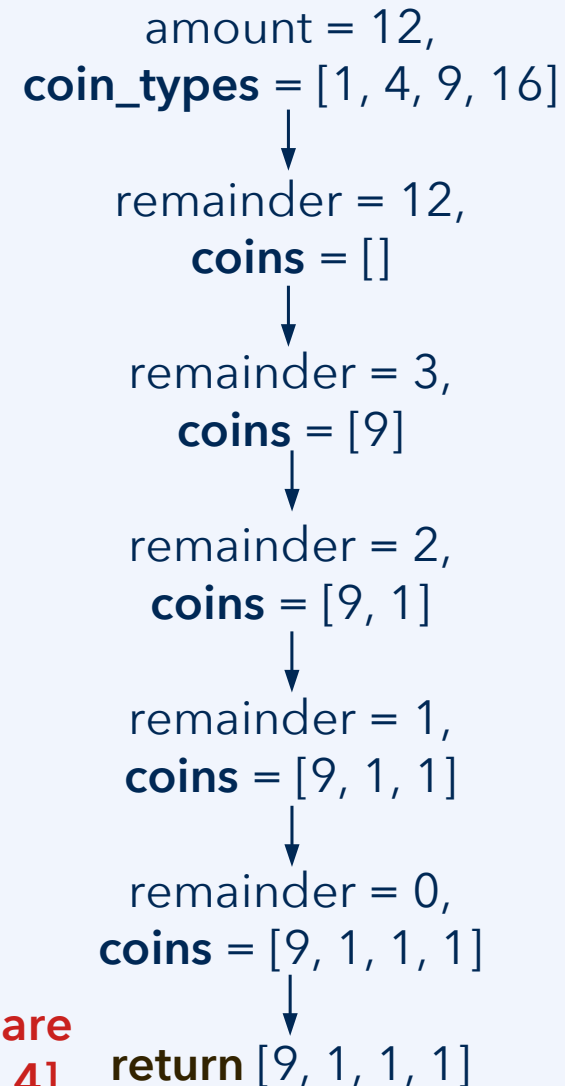
Cashier problem (T2.2)

Condition for the greedy algorithm:

The shortest sum containing coin x never consists of more coins than the shortest equivalent sum containing only coins $< x$.

```
def cashier(amount, coin_types):
    coins = []
    remainder = amount
    while remainder >= coin_types[0]:
        for i in range(len(coin_types)-1, -1, -1):
            if remainder >= coin_types[i]:
                coins.append(coin_types[i])
                remainder -= coin_types[i]
            break
    return coins
```

compare
[4, 4, 4]



Cashier problem (T2.2)

greedy algorithm

Let n = amount and k = $\text{len}(\text{coin_types})$.

```
def cashier(amount, coin_types):
```

```
    coins = []
```

$O(1)$ instructions

```
    remainder = amount
```

```
    while remainder >= coin_types[0]:
```

Upper bound: $O(n)$ iterations

```
        for i in range(len(coin_types)-1, -1, -1):
```

– Upper bound: $O(k)$ iterations

```
            if remainder >= coin_types[i]:
```

```
                coins.append(coin_types[i])
```

```
                remainder -= coin_types[i]
```

```
                break
```

```
    return coins
```

- $O(1)$ time on average, for a well-managed **dyn. array**.
 $O(1)$ time worst case if we were using a **linked list**.

Cashier problem (T2.2)

greedy algorithm

Let $n = \text{amount}$ and $k = \text{len}(\text{coin_types})$.

```
def cashier(amount, coin_types):
```

```
    coins = []
```

$O(1)$ instructions

```
    remainder = amount
```

```
    while remainder >= coin_types[0]:
```

Upper bound: $O(n)$ iterations

```
        for i in range(len(coin_types)-1, -1, -1):
```

– Upper bound: $O(k)$ iterations

```
            if remainder >= coin_types[i]:
```

```
                coins.append(coin_types[i])
```

```
                remainder -= coin_types[i]
```

```
                break
```

```
    return coins
```

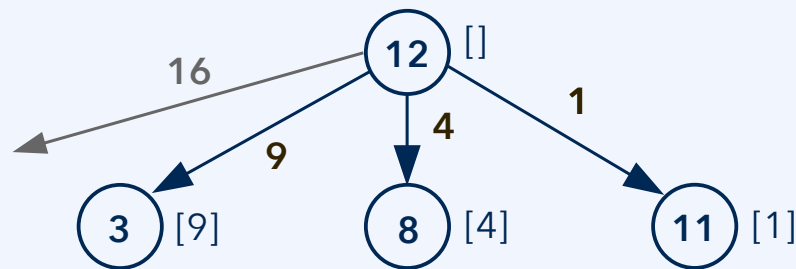
- $O(1)$ time* per iteration
[*rigorous with a linked list]

$O(kn)$ time efficiency

that is $O(n)$ if k is treated as constant

Cashier problem (T2.2)

Illustration: Dynamic programming algorithm for the cashier problem



amount = 12,
coin_types = [1, 4, 9, 16]

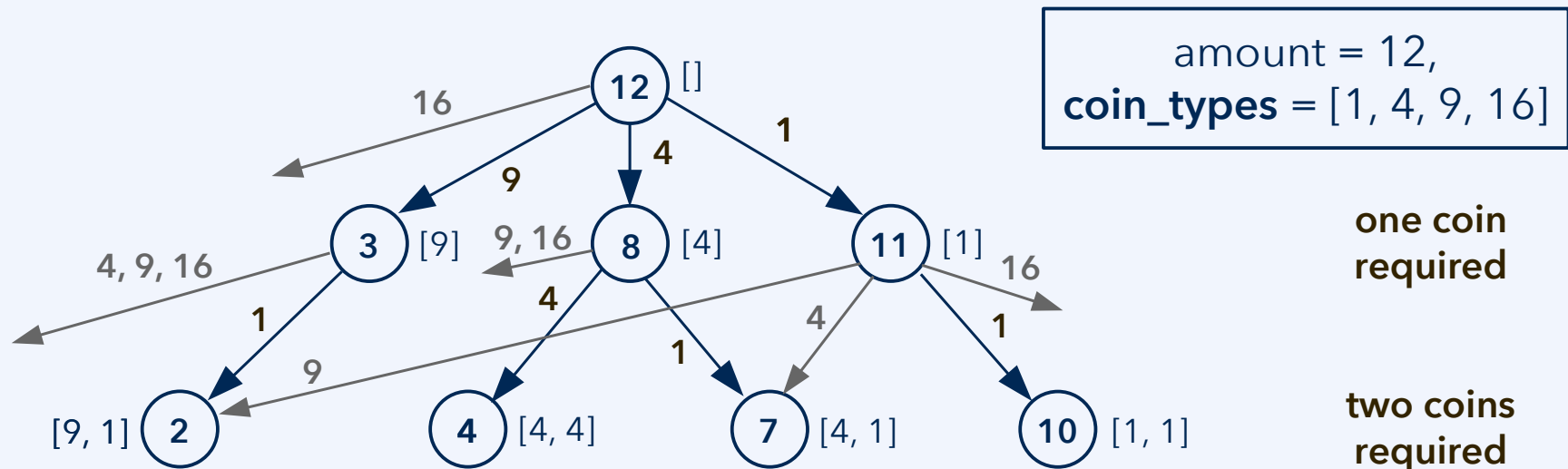
If we return a 9-valued coin, the remainder reduces to 3.

If we return a 4-valued coin, the remainder reduces to 8.

If we return a 1-valued coin, the remainder reduces to 11.

Cashier problem (T2.2)

Illustration: Dynamic programming algorithm for the cashier problem



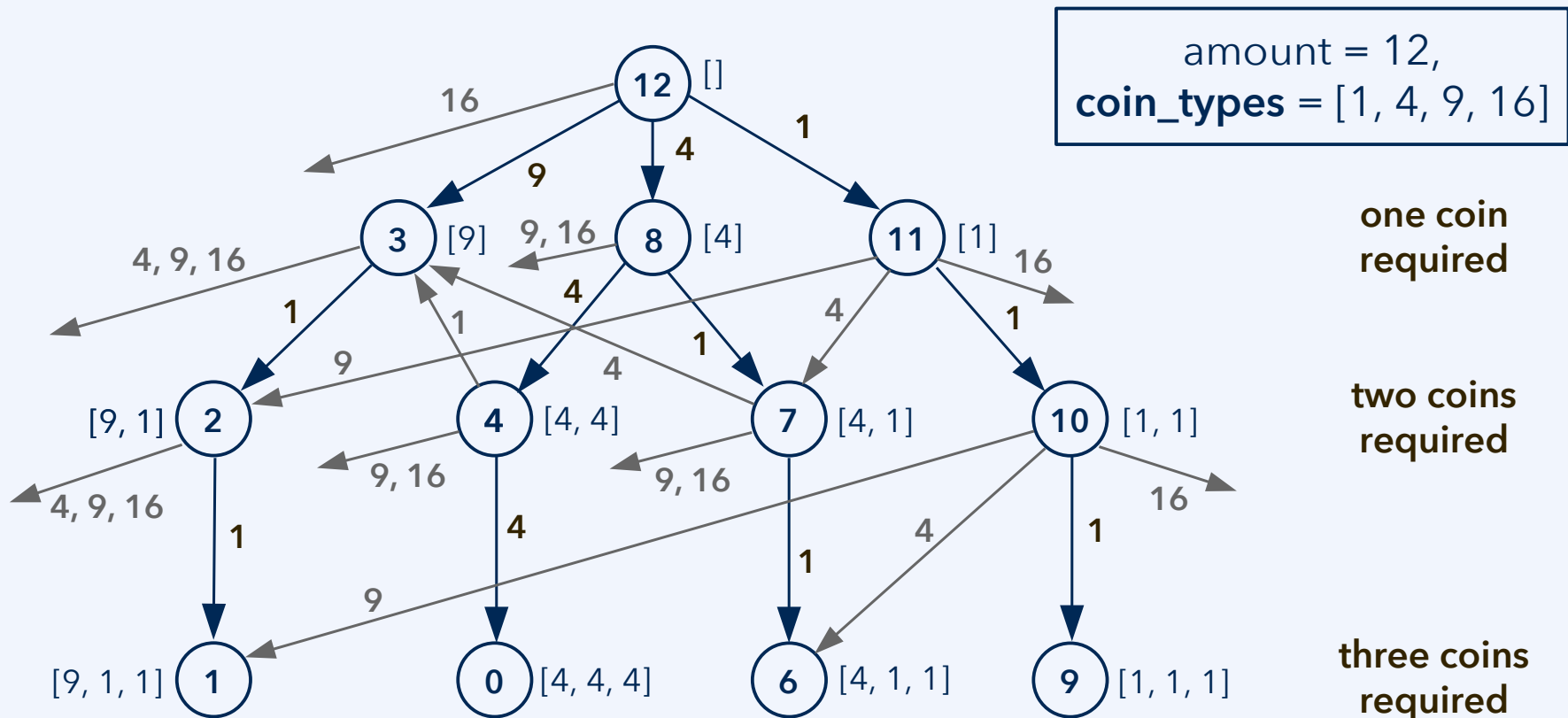
A remainder of 12 currency units can be reached using zero coins.

A remainder of 3, 8, or 11 can be reached using one coin.

A remainder of 2, 4, 7, or 10 can be reached using two coins.

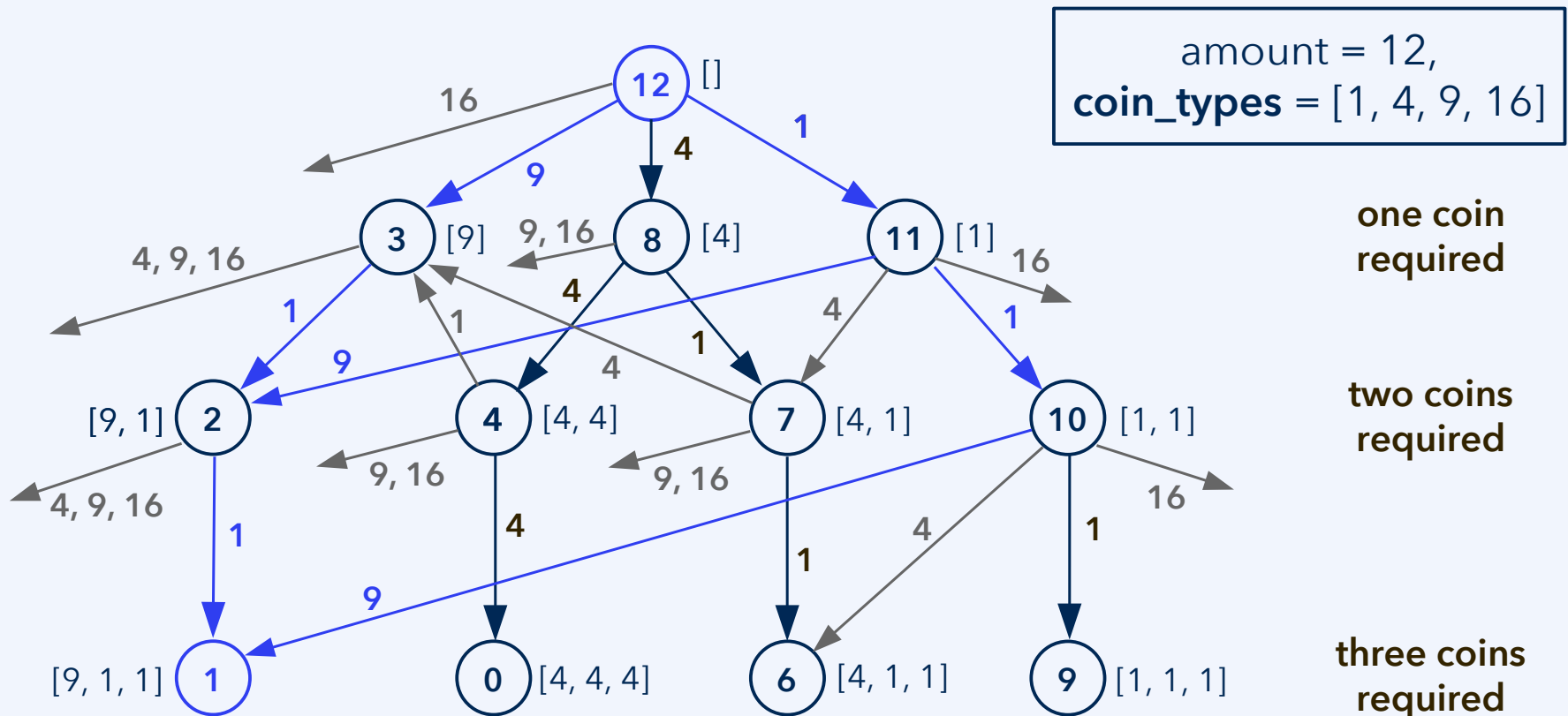
Cashier problem (T2.2)

Illustration: Dynamic programming algorithm for the cashier problem



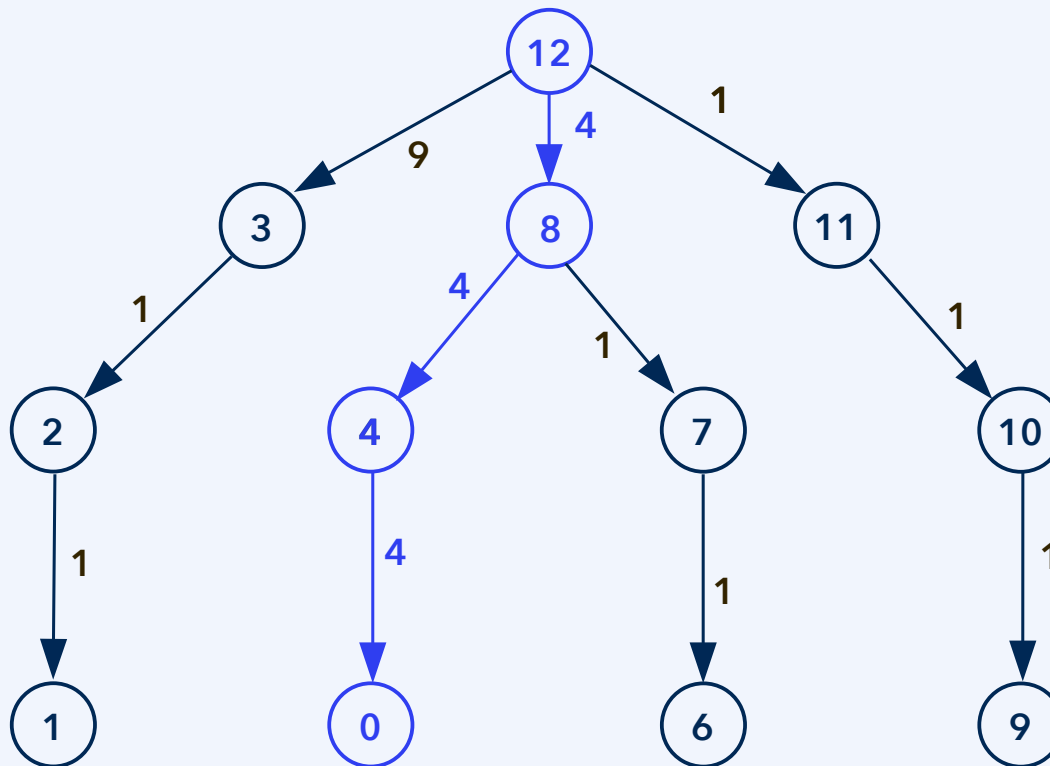
Cashier problem (T2.2)

Overlapping subproblems, e.g., equivalence of $[9, 1, 1]$, $[1, 9, 1]$, and $[1, 1, 9]$



Cashier problem (T2.2)

Illustration: Dynamic programming algorithm for the cashier problem



This reduces to computing a **BFS spanning tree**, over a graph with at most $n+1$ nodes, with node out-degree upper bound of k .

Time efficiency $O(kn)$, same as for the greedy algorithm.