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## INF205

## **Resource-efficient programming**

- 3 Concurrency
- 3.1 Parallel programming
- 3.2 Message passing interface (MPI)

## **Concurrency: Why does it matter?**



Assume a scenario where we can split a code into a fraction f that can be parellelized and the remainder 1 - f that is always sequential, never parallel.

Adding two vectors c[i] = a[i] + b[i], for *i* from 0 to 99 999, can be parallelized. Waiting for new instructions from the user cannot be parallelized.

**Speedup** is the *factor by which runtime decreases*; here, due to parallelization.

### Amdahl's law:

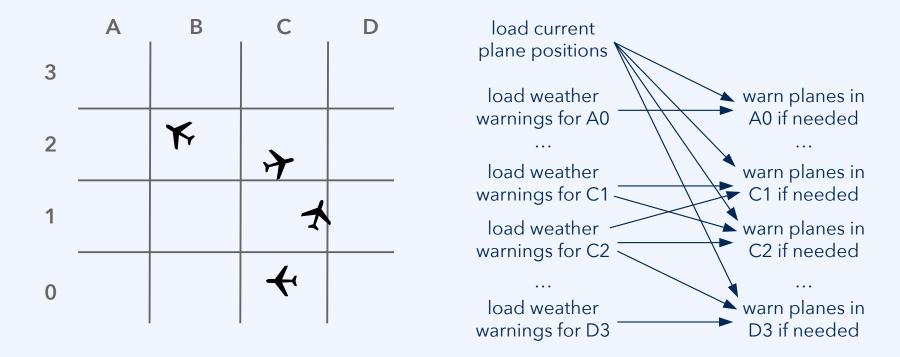
- Runtime with a single process is given by some  $t_1 = (1-f)t_1 + ft_1$ .
- Now assume that we are parallelizing the code as perfectly as possible:
  - With *n* parallel processes, the runtime becomes  $t_n = (1-f)t_1 + \frac{ft_1}{n}$ .
- Now assume that we have infinite computing resources at our hands:
  - With infinite parallel processes, the runtime becomes  $t_{m} = (1 f) t_{1}$ .
- The maximum possible speedup for our code is  $S_{\infty} = t_{\infty}/t_1 = 1/(1-f)$ .

If f = 99% can be parallelized, speedup can never be greater than  $S_{m} = 100$ .

## **Concurrency: Why does it matter?**

Parallel programming is about efficiently exploiting a parallel architecture.

Concurrency as a topic is about what can be parallelized and what cannot.

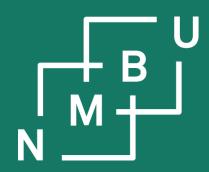


Domain decomposition is one of the techniques for this kind of concurrency.

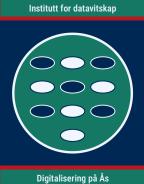
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3 Concurrency

## 3.1 Parallel programming





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### Message passing

Message passing is the most general paradigm of parallel programming.

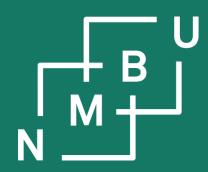
It can be carried out *irrespective whether* or not the *processes* (can also be called **ranks** in MPI) are executed on the same computing node and *have shared memory access*. It only assumes that they can exchange messages.

Challenges of message passing based parallelization:

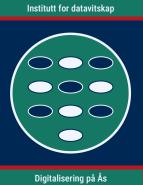
- Synchronization (waiting) while processes need to talk to each other.
- What if there are very many processes, do they all message each other?
- What if the recipient would already have had access to the data?
- Processes need to figure out what information they must give to others.

In high performance computing, message-passing based parallelization is usually done using **MPI**, the message passing interface.

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## 3 Concurrency

# 3.1 Parallel programming3.2 Message passing interface





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### **MPI: Getting started**

The target systems of MPI programs are often *clusters with thousands of cores*.

However, the code is not usually developed on these systems, but on the programmers' usual working environment. Even on a laptop/workstation, MPI makes you realize a *speedup*, since today these are all *multicore systems*.

To get started install an MPI environment, e.g., **Open MPI** (package **openmpi**).

The **compiler command** becomes "**mpiCC** ..." or similar (instead of "g++ ..."). The *binary executable* produced by the compiler *will not run on its own*!

Instead: mpirun -np <number of processes> <executable>

This creates a number of parallel processes with ranks starting from 0. Often the *process with rank 0* takes the role of the "master" or "scheduler".

See also the Open MPI documentation: https://www.open-mpi.org/doc/v4.1/

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## MPI: Getting started

An MPI program needs to *initialize* and *finalize* the MPI environment. Every process needs to *know its rank* (and, usually, the *number of processes*).

```
#include <mpi.h>
int main(int argc, char** argv)
{
    MPI_Init(&argc, &argv);
    int rank = 0; // what is the rank of this process?
    MPI_Comm_rank(MPI_COMM_WORLD, &rank);
    int size = 0; // how many processes are there?
    MPI_Comm_size(MPI_COMM_WORLD, &size);
    ... // here comes the actual program
    MPI_Finalize();
```

(See the **mpi-primes** example code.)

Often the rank no. of a process, together with the number of processes, is already enough input to implement a basic parallelization scheme.

This is also the case for our prime-number test example:

5 7 11 13 17 19 23... 0 0 1 1 2 2 3...





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### **MPI send and receive**

The most basic communication step is send/receive from one rank to another.

### int MPI\_Send(

void\* content, int count, MPI\_Datatype type, int destination\_rank, int tag, MPI\_Comm handle );

**content** is the address from which the source data are read; it is often an array, but can also be a pointer to a single data item

### int MPI\_Recv(

void\* buffer, int count, MPI\_Datatype type, int source\_rank, int tag, MPI\_Comm handle, MPI\_Status\* status );

**buffer** is an address to which the received data can be written; the programmer needs to take care of memory allocation, *etc*.

#### count is the number of data items

**type** is their type as an MPI environment expression (e.g., MPI\_SHORT\_INT, MPI\_INT64\_T, MPI\_FLOAT, ...)

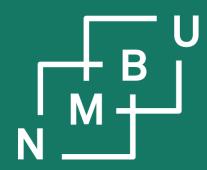
tag is an identifier; send and receive must have the same tag

**destination\_rank** is the rank of the process with the matching MPI\_Recv(...) operation

source\_rank is the rank of the process with
the matching MPI\_Send(...) operation

(Standard values from handle and status are MPI\_COMM\_WORLD and MPI\_STATUS\_IGNORE.)

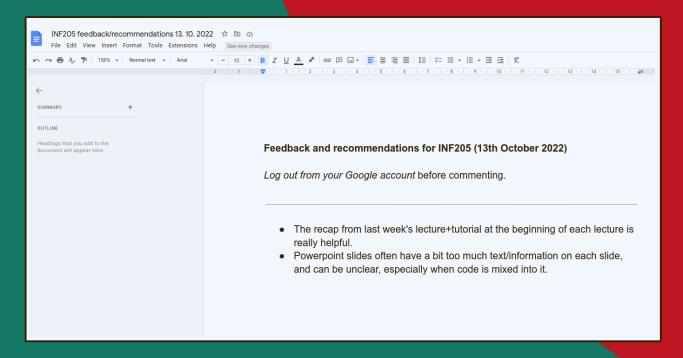
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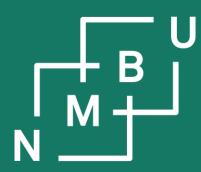
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## Conclusion



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